Data Types with Symmetries via Action Containers

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Overview

Goal of the talk

Introduce action containers to model data types with symmetries

Contents

- Background
 - Endofunctors and algebraic data types
 - Containers for polynomial functors
- Action containers
 - Construction via universal property
 - Closure properties
- 2-categorical interpretation: Embedding as 2-endofunctors of groupoids

Endofunctors model algebraic data types

Many data types can be modeled in the category of Set-endofunctors:

$$\operatorname{List}(X) := \sum_{n:\mathbb{N}} (\operatorname{Fin}(n) \to X) \quad \operatorname{Maybe}(X) := 1 + X \quad \operatorname{RoseTree}(X) := \mu Y. X + \operatorname{List}(Y)$$

New endofunctors can be built by "algebraic" manipulations:

"a pair of an
$$F$$
 and a G " "either an F or a G " "a K -tuple of F s"
 $F \times G$ $F + G$ F^K

Containers: syntax for polynomials

The nice class of *polynomial* endofunctors is described by *containers*:



Sanity Check

 $\llbracket - \rrbracket$: **Container** \rightarrow Endo(**Set**) is a fully faithful functor.

Morphisms of containers describe exactly the morphisms of their interpretations.

Non-polynomial endofunctors

Caveat

Not all interesting functors are covered by this framework.

Example

Finite multisets are not polynomial:

$$\mathsf{FMSet}(X) := \sum_{n:\mathbb{N}} (\mathsf{Fin}(n) \to X) / \sim$$

The relation is generated by identifying tuples up to permutation:

$$(x_1,\ldots,x_n)\sim (x_{\pi 1},\ldots,x_{\pi n})\quad \forall \pi:\mathsf{Fin}(n)\simeq\mathsf{Fin}(n)$$

Action containers

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Definition
An action container F = (S \triangleright P \triangleleft^{\sigma} G) consists of
shapes a set S
positions a family of sets P : S \rightarrow \mathbf{Set}
symmetries a family of groups G : S \rightarrow \mathbf{Group}
actions a family of group actions: for each s : S, \sigma_s is an action of G_s on P_s
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Intuition

Symmetries tell us under which permutations of positions the data type is invariant.

Interpretation

$$\llbracket S \triangleright P \triangleleft^{\sigma} G \rrbracket(X) := \sum_{s:S} (P_s \to X) / \sim_s \qquad v \sim_s w := \exists g : G_s. v = w \circ \sigma_s(g)$$

Group actions?

Definition

An action of a group G on a set X is a group homomorphism $\sigma : G \to \mathfrak{S}(X)$, where $\mathfrak{S}(X)$ is the group of *automorphisms* $X \simeq X$.

Equivalently, a G-action on X is a functor $\mathbb{B}\sigma:\mathbb{B}G
ightarrow$ Set

- $\mathbb{B}G$ is G seen as a 1-object groupoid
- the lone object is sent to X
- ▶ loops $g : \bullet \xrightarrow{\sim} \bullet$ are sent to $\sigma(g) : X \xrightarrow{\sim} X$

Example: Finite multisets

Finite multisets are "lists up to permutation", thus come from *the* permutation action:

$$\mathsf{FMSet} = (n: \mathbb{N} \triangleright \mathsf{Fin}(n) \triangleleft^{\mathsf{id}} \mathfrak{S}(\mathsf{Fin}(n))) \qquad \mathsf{id}: \mathfrak{S}(\mathsf{Fin}(n))
ightarrow \mathfrak{S}(\mathsf{Fin}(n))$$

Its interpretation:

$$\llbracket \mathsf{FMSet}
rbracket X = \sum_{n:\mathbb{N}} X^n / {\sim}$$

where (\sim) is generated by

$$(x_1,\ldots,x_n)\sim (x_{\pi(1)},\ldots,x_{\pi(n)})\quad \forall\pi:\mathsf{Fin}(n)\simeq\mathsf{Fin}(n)$$

Example: Cyclic lists

Cyclic lists come from a $\mathbb Z\text{-}\mathsf{action}$ on finite sets:

 $\mathsf{Cyc} = (n : \mathbb{N} \triangleright \mathsf{Fin}(n) \triangleleft^{\sigma_n} \mathbb{Z})$ $\sigma_n : \mathbb{Z} \to \mathfrak{S}(\mathsf{Fin}(n))$

where σ_n is generated from the successor automorphism,

$$\operatorname{suc}_n : \operatorname{Fin}(n) \simeq \operatorname{Fin}(n)$$

 $\operatorname{suc}_n(x) := x + 1 \mod n$
 $\sigma_n(k) := \underbrace{\operatorname{suc}_n \circ \cdots \circ \operatorname{suc}_n}_{k \text{ times}}$

In its interpretation,

$$\llbracket \mathsf{Cyc} \rrbracket X = \sum_{n:\mathbb{N}} X^n / \sim,$$

the relation (\sim) is generated by

$$(x_1,\ldots,x_n)\sim (x_n,x_1,\ldots,x_{n-1})$$

Morphisms of action containers

A morphism of action containers preserves shapes, positions, and symmetries:

Definition

Let $F = (S \triangleright P \triangleleft^{\sigma} G)$, $G = (T \triangleright Q \triangleleft^{\tau} H)$. A morphism $(u \triangleright f \triangleleft \varphi) : F \rightarrow G$ consists of

- ▶ a function on shapes $u: S \to T$
- ▶ a family of functions $f : \prod_{s:S} Q_{us} \to P_s$
- ▶ a family of group homomorphisms $\varphi : \prod_{s:S} G_s \rightarrow H_{us}$

such that the following diagram commutes for all s : S, $g : G_s$:

$$\begin{array}{ccc} Q_{us} & \stackrel{f_s}{\longrightarrow} & P_s \\ \tau_{us}(\varphi_s(g)) & & & \downarrow \sigma_s(g) \\ Q_{us} & \stackrel{f_s}{\longrightarrow} & P_s \end{array}$$

Our definitions are inspired by *quotient containers*,¹ but there are differences:

Differences on objects

In a quotient container $(S \triangleright P/G)$, the symmetries are restricted to subgroups of permutation groups:

$$\iota_s: G_s \leq \mathfrak{S}(P_s)$$

Action containers can describe symmetry groups *larger* than $\mathfrak{S}(P_s)$ (e.g. the \mathbb{Z} -action in Cyc).

Comparision to quotient containers II

Differences on morphisms

Morphisms of quotient containers do not respect symmetries: for all g : G_s there merely exists some h : H_{us} such that

$$\begin{array}{c} Q_{us} \xrightarrow{f_s} P_s \\ \downarrow_{us}(h) \downarrow \qquad \qquad \qquad \downarrow_{\iota_s}(g) \\ Q_{us} \xrightarrow{f_s} P_s \end{array}$$

as opposed to the output of a group homomorphism $h \doteq \varphi(g)$.

Morphisms of quotient containers are *quotiented* by some relation, such that [[−]]_/: QuotCont → Endo (Set) is fully faithful.

¹Abbott, Altenkirch, Ghani, and McBride, "Constructing Polymorphic Programs with Quotient Types".

Construction from universal property

The category of action containers is "just" families of group actions.

- For each shape s : S, there is some group (G_s) acting (σ_s) on some set (P_s) .
- ln a morphism, for each shape s there is some homomorphism (φ_s) and a function (f_s) that respect each other.

This suggests a modular construction.

A category of group actions

We define a category Action of group actions and equivariant maps:

- Objects are triples (G, P, σ) , where σ is a *G*-action on *P*
- ▶ A morphism $(\varphi, f) : (G, P, \sigma) \rightarrow (H, Q, \tau)$ consists of
 - ▶ a group homomorphism φ : $G \rightarrow H$,
 - ▶ a function $f : P \leftarrow Q$,

such that f is equivariant, i.e. for all g : G

$$\begin{array}{ccc} Q & \stackrel{f}{\longrightarrow} P \\ & & \downarrow \\ T(\varphi(g)) \downarrow & & \downarrow \\ Q & \stackrel{f}{\longrightarrow} P \end{array}$$

► Morphisms compose by pasting equivariance squares. This is like the category of *G*-sets, except *G* can vary.

Universal property

Definition (Fam-construction)

```
The free coproduct completion Fam(C) of a category C:
objects pairs (S, x) of a set S and a family x : S \to ob(C)
morphisms for (u, f) : (S, x) \to (T, y), a function u : S \to T and a family of
morphisms f : \prod_{s:S} C(x_s, y_{us})
```

Theorem

The category of action containers is equivalent to the free coproduct completion of the category of group actions:

$\textbf{ActionCont} \simeq \mathsf{Fam}(\textbf{Action})$

Presenting ActionCont as Fam(Action) gives us closure properties:

Proposition

Action containers are closed under (arbitrary) coproducts and products.

Proof.

- coproducts: by construction
- products: Action is closed under products

Closure under exponentiation

Definition

The constant container at a set J is $\mathbf{k}J := (J \triangleright 0 \triangleleft^{\mathsf{id}} 1)$

Why constant? Answer: $\llbracket \mathbf{k} J \rrbracket(X) \simeq J$.

Proposition

Action containers are closed under exponentiation by constants: For any container F and set J, there is a container F^J and a universal morphisms eval : $\mathbf{k}J \times F^J \to F$.

A model of strictly positive types

Action containers model non-inductive single-variable strictly positive types.²

- strictly positive types are closed under products F × G, coproducts F + G and constant exponentiation F^J.
- ▶ single-variable: we do not consider *indexed* container³
- ► non-inductive: have not yet looked at smallest µX. F(X, -) and largest νX. F(X, -) fixedpoints yet.

²Abbott, Altenkirch, and Ghani, "Containers: Constructing strictly positive types".

³Conceptually not too hard, but requires more bookkeeping

Properties of the interpretation

Caveat

The interpretation functor is no longer fully faithful.

Reason

Quotients in **Set** forget why morphisms have been identified. The evidence is an element of a symmetry group, G_s .

Fix

Make the evidence part of the data, and interpret action containers in 2-endofunctors of groupoids.

Interpretation in groupoids

To interpret action containers in groupoids:

- 1. Enhance ActionCont into a 2-category, in small steps
- 2. Embed them into the 2-category of symmetric containers⁴
- 3. Compose with interpretation of the latter:

$$\textbf{ActionCont} \longrightarrow \textbf{SymmCont} \xleftarrow{\llbracket - \rrbracket} \mathsf{Endo}(\mathsf{hGpd})$$

⁴Gylterud, "Symmetric Containers".

A 2-category of action containers

Present the 2-category ActionCont again as families of group actions:

Group forms a 2-category⁵ with morphisms "up to conjugation":

$$\operatorname{\mathsf{Group}}_2(arphi,\psi):=\sum_{r:H}arphi=r\psi r^{-1}\quad \forall arphi,\psi:\operatorname{\mathsf{Group}}_1(G,H)$$

(*r* is called a *conjugator*)

Similarly for Action: 2-cells are conjugators that relate equivariant maps.

▶ 2-cells of Fam *C* are families of 2-cells of *C*.

We define

ActionCont := Fam(Action)

⁵Hofstra and Karvonen, "Inner automorphisms as 2-cells".

Symmetric containers in HoTT

Definition A symmetric container $(S \triangleleft P)$ consists of shapes an *h*-groupoid⁶ S positions a function $P : S \rightarrow hSet$

In HoTT, symmetries are internalized as paths:

• paths s = t in S encode symmetries on shapes

symmetries of positions are induced functorially:

$$\operatorname{cong}(P): s = t
ightarrow P(s) = P(t)$$

interpretation in h-groupoids:

$$\llbracket S \triangleleft P \rrbracket X := \sum_{s:S} P(s) \to X$$

⁶at most one proof p = q for all p, q : s = t in S

The Delooping-construction

A group G defines a 1-object h-groupoid $\mathbb{B}G$, implemented as a higher inductive type:

• :
$$\mathbb{B}G$$
 $g: G$ $g, h: G$ $g, h: G$ $g, h: G$ $g, h: G$

By recursion on $\mathbb{B}G$, each action $\sigma: G \rightarrow \mathfrak{S}(X)$ defines a family

 $\overline{\mathbb{B}}\sigma:\mathbb{B}G\to\mathsf{hSet}$

Theorem

The above extend to 2-functors

 $\mathbb{B}: \textbf{Group} \to \mathsf{hGpd}$

 $\overline{\mathbb{B}}$: Action \rightarrow SymmCont $\overline{\mathbb{B}}(G, \sigma) := (\mathbb{B}G \triangleleft \overline{\mathbb{B}}\sigma)$

Both are locally weak equivalences.

Symmetric containers from action containers I

We can embed single actions as symmetric containers:

Action
$$\overset{\bar{\mathbb{B}}}{\longrightarrow}$$
 SymmCont

We can lift this to families of such (i.e. action containers):

Theorem

The lifting $\mathsf{Fam}(\bar{\mathbb{B}})$: ActionCont \rightarrow $\mathsf{Fam}(SymmCont)$ is

1. locally fully faithful

2. locally a weak equivalence, assuming the axiom of choice

Symmetric containers from action containers II

What's missing?

$$\textbf{ActionCont} \ \ \overset{\bar{\mathbb{B}}}{\longleftarrow} \ \ \mathsf{Fam}(\textbf{SymmCont}) \ \ \overset{???}{\longrightarrow} \ \ \textbf{SymmCont}$$

Proposition

There is a 2-functor, summing families of symmetric containers:

 $\Sigma:\mathsf{Fam}(\mathsf{SymmCont})\to\mathsf{SymmCont}$

The local functors

$$\Sigma_1: \mathsf{Fam}_1(X, Y) o \mathbf{SymmCont}_1(\Sigma_0 X, \Sigma_0 Y)$$

are fully faithful if all shape groupoids of X are connected.

Interpretation in groupoids

We can *locally* embed action containers in 2-endofunctors:

Theorem The factorization

$$\textbf{ActionCont} \xrightarrow{\text{Fam}(\tilde{\mathbb{B}})} \text{Fam}(\textbf{SymmCont}) \xrightarrow{\Sigma} \textbf{SymmCont} \xrightarrow{\mathbb{I}-\mathbb{I}} \text{Endo}(h\text{Gpd})$$

is locally fully faithful.

Proof.

In the image of $Fam(\overline{\mathbb{B}})$, shape groupoids are connected.

This fully classifies 1- and 2-cells of action containers.

Conclusion

We have...

- constructed action containers via a universal property
- showed closure under desirable operations
- connected them to symmetric containers
- embedded them as 2-endofunctors
- developed all of this in Cubical Agda:
 - includes formalization of necessary 2-category theory

For a draft and the formalization:



https://phijor.me/publications/

 ${\tt 2025-data-types-with-symmetries-via-action-containers.html}$

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